

Some designs and binary codes preserved by the simple group R_u of Rudvalis

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Motivation

- The simple group Ru of Rudvalis is one the 26 sporadic simple groups.
- It has a rank-3 primitive permutation representation of degree 4060 which can be used to construct a strongly regular graph Γ with parameters $v = 4060$, $k = 1755$, $\lambda = 730$ and $\mu = 780$ or its complement a strongly regular $\tilde{\Gamma} = (4060, 2304, 1328, 1280)$ graph.
- The stabilizer of a vertex u in this representation is a maximal subgroup isomorphic to the Ree group $2F_4(2)$ producing orbits $\{u\}$, Δ_1 , Δ_2 of lengths 1, 1755, and 2304 respectively. The regular graphs Γ , $\tilde{\Gamma}$, Γ^R , $\tilde{\Gamma}^R$, Γ^S are constructed from the sets Δ_1 , Δ_2 , $\{u\} \cup \Delta_1$, $\{u\} \cup \Delta_2$, and $\Delta_1 \cup \Delta_2$, respectively.

Motivation

- If A denotes an adjacency matrix for Γ then $B = J - I - A$, where J is the all-one and I the identity 4060×4060 matrix, will be an adjacency matrix for the graph $\tilde{\Gamma}$ on the same vertices.
- We examine the neighbourhood designs \mathcal{D}_{1755} , \mathcal{D}_{1756} , \mathcal{D}_{2304} , \mathcal{D}_{2305} and \mathcal{D}_{4059} and corresponding binary codes \mathcal{C}_{1755} , \mathcal{C}_{1756} , \mathcal{C}_{2304} , \mathcal{C}_{2305} , and \mathcal{C}_{4059} defined by the binary row span of A , $A + I$, B , $B + I$ and $A + B$ respectively.

Background - t -(v, k, λ) Designs

- An incidence structure $\mathcal{D} = (\mathcal{P}, \mathcal{B}, \mathcal{I})$ with **point set** \mathcal{P} and **block set** \mathcal{B} and incidence $\mathcal{I} \subseteq \mathcal{P} \times \mathcal{B}$ is a t - (v, k, λ) design if
 - $|\mathcal{P}| = v$;
 - every block $B \in \mathcal{B}$ is incident with precisely k points;
 - every t distinct points are together incident with precisely λ blocks. t, v, k and λ are non-negative integers;
 - $|\mathcal{B}| = b$ is the number of blocks;

An **incidence matrix** for \mathcal{D} is a $b \times v$ matrix $A = (a_{ij})$ of 0's and 1's such that

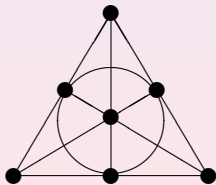
$$a_{ij} = \begin{cases} 1 & \text{if } (p_j, B_i) \in \mathcal{I} \\ 0 & \text{if } (p_j, B_i) \notin \mathcal{I} . \end{cases}$$

The Fano Plane is a $2 - (7, 3, 1)$ Design

Take $S = \{1, 2, 3, 4, 5, 6, 7\}$ and consider the subsets:
 $\{1, 2, 4\}$ $\{2, 3, 5\}$ $\{3, 4, 6\}$, $\{4, 5, 7\}$ $\{5, 6, 1\}$ $\{6, 7, 2\}$ $\{7, 1, 3\}$.

We have a $2 - (7, 7, 3, 3, 1)$ -design. We can have a geometrical interpretation of this design as follows:

- The elements of $1, 2, 3, \dots, 7$ are represented by points and the blocks by lines (6 straight lines and a circle). This is known as the projective plane of order 2.



Incidence matrix - an example

		Blocks (lines)						
		b_1	b_2	b_3	b_4	b_5	b_6	b_7
Points	p_1	1	0	0	0	1	0	1
	p_2	1	1	0	0	0	1	0
	p_3	0	1	1	0	0	0	1
	p_4	1	0	1	1	0	0	0
	p_5	0	1	0	1	1	0	0
	p_6	0	0	1	0	1	1	0
	p_7	0	0	0	1	0	1	1

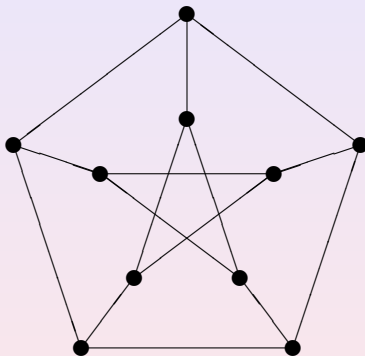
Table : Incidence matrix of the 2 – (7, 3, 1) Design

Background - Graphs

- A **graph** $\mathcal{G} = (V, E)$, consists of a finite set of vertices V together with a set of edges E , where an edge is a subset of the vertex set of cardinality 2. Our graphs are undirected.
- The **valency** of a vertex is the number of edges containing the vertex.
- A graph is **regular** if all the vertices have the same valency; a regular graph is **strongly regular** of type (n, k, λ, μ) if it has n vertices, valency k , and if any two adjacent vertices are together adjacent to λ vertices, while any two non-adjacent vertices are together adjacent to μ vertices.
- The **adjacency matrix** $A(\mathcal{G})$ of \mathcal{G} is the $n \times n$ matrix with

$$(i, j) = \begin{cases} 1 & \text{if } x_i \text{ and } x_j \text{ are adjacent,} \\ 0 & \text{otherwise.} \end{cases}$$

The Petersen Graph is strongly regular



Error-correcting codes

- Let F be any set of size q and let F^n denote the set of n -tuples of elements of F (usually here F will be a finite field). Call the elements of F^n vectors.
- A **q -ary code** C of length n is a set of elements of F^n , called **codewords** or vectors, and written $x_1x_2 \dots x_n$, or (x_1, x_2, \dots, x_n) , where $x_i \in F$ for $i = 1, \dots, n$.

Definition

Let $v = (v_1, v_2, \dots, v_n)$ and $w = (w_1, w_2, \dots, w_n)$ be in F^n . The **Hamming distance**, $d(v, w)$, between v and w is the number of coordinate places in which they differ:

$$d(v, w) = |\{i \mid v_i \neq w_i\}|.$$



Error Correcting Codes

Definition

The *minimum distance* $d(C)$ of a code C is the smallest of the distances between distinct codewords; i.e.

$$d(C) = \min\{d(v, w) \mid v, w \in C, v \neq w\}.$$

Theorem

If $d(C) = d$ then C can detect up to $d - 1$ errors or correct up to $\lfloor (d - 1)/2 \rfloor$ errors.

Linear Codes

- A code C over the finite field $F = \mathbf{F}_q$ of order q , of length n is **linear** if C is a subspace of $V = F^n$. If $\dim(C) = k$ and $d(C) = d$, then we write $[n, k, d]$ or $[n, k, d]_q$ for the q -ary code C .
- If C is a q -ary $[n, k]$ code, a **generator matrix** for C is a $k \times n$ array obtained from any k linearly independent vectors of C .
- Let C be a q -ary $[n, k]$ code. The **dual code** of C is denoted by C^\perp and is given by

$$C^\perp = \{v \in F^n \mid (v, c) = 0 \text{ for all } c \in C\}.$$

Linear Codes- Continued

- A **check matrix** for C is a generator matrix H for C^\perp .
- Two linear codes of the same length and over the same field are **isomorphic** if they can be obtained from one another by permuting the coordinate positions.
- An **automorphism** of a code C is an isomorphism from C to C .
- Any code is isomorphic to a code with generator matrix in **standard form**, i. e. the form $[I_k | A]$; a check matrix then is given by $[-A^T | I_{n-k}]$. The first k coordinates are the **information symbols** and the last $n - k$ coordinates are the **check symbols**.



A preliminary result

Result

Let G be a finite primitive permutation group acting on the set Ω of size n . Let $\alpha \in \Omega$, and let $\Delta \neq \{\alpha\}$ be an orbit of the stabilizer G_α of α . If $\mathcal{B} = \{\Delta^g \mid g \in G\}$ and, given $\delta \in \Delta$, $\mathcal{E} = \{\{\alpha, \delta\}^g \mid g \in G\}$, then $\mathcal{D} = (\Omega, \mathcal{B})$ forms a symmetric 1 - $(n, |\Delta|, |\Delta|)$ design. Further, if Δ is a self-paired orbit of G_α then $\Gamma = (\Omega, \mathcal{E})$ is a regular connected graph of valency $|\Delta|$, \mathcal{D} is self-dual, and G acts as an automorphism group on each of these structures, primitive on vertices of the graph, and on points and blocks of the design.

In fact one can use any union of orbits of a point-stabilizer in this construction, and this is the approach that we will adopt in the paper.



The Rudvalis group Ru

The primitive representations Ru are listed in Table 2. The first column gives the ordering of the primitive representations; the second gives the maximal subgroups; the third gives the degree (the number of cosets of the point stabilizer);

No.	Max. sub.	Deg.	No.	Max. sub.	Deg.
1	$2_{F_4(2)}$	4060	9	$L_2(29)$	11980800
2	$(2^6:U_3 3):2$	188500	10	$5^2:4S_5$	12160512
3	$(2^2 \times S_Z(8)):3$	417600	11	$3 \cdot A_6 \cdot 2^2$	33779200
4	$2^{3+8}:L_3(2)$	424125	12	$5_+^{1+2}:[2^5]$	36481536
5	$U_3(5):2$	579072	13	$L_2(13):2$	66816000
6	$2 \cdot 2^{4+6}:S_5$	593775	14	$A_6 \cdot 2^2$	101337600
7	$L_2(25) \cdot 2^2$	4677120	15	$5:4 \times A_5$	121605120
8	A_8	7238400			

Table : Maximal subgroups of Ru

The graphs, designs and codes

- The above Table shows that there is just one class of maximal subgroups of R_u of index 4060. The stabilizer of a vertex u in this representation is a maximal subgroup isomorphic to $2_{F_4(2)}$, producing orbits $\{u\}$, Δ_1 , and Δ_2 of lengths 1, 1755 and 2304 respectively.
- The regular graphs Γ , Γ^R , $\tilde{\Gamma}$, $\tilde{\Gamma}^R$ are constructed from the sets Δ_1 , $\{u\} \cup \Delta_1$, Δ_2 and $\{u\} \cup \Delta_2$, respectively.
- The binary codes C_{1755} , C_{1756} , C_{2304} , C_{2305} whose properties we will be examining are obtained as described below.

The graphs, designs and codes

- The rows of an adjacency matrix A for Γ give the blocks of the neighbourhood design of Γ which we will denote \mathcal{D}_{1755} . Notice that \mathcal{D}_{1755} is a self-dual symmetric 1-(4060, 1755, 1755) design. We write \mathcal{C}_{1755} to denote the binary code spanned by the rows of \mathcal{D}_{1755} .
- From the rows of an adjacency matrix $A + I$ of the reflexive graph Γ^R we obtain the self-dual symmetric 1-(4060, 1756, 1756) design \mathcal{D}_{1756} , and the binary code \mathcal{C}_{1756} .
- The rows of an adjacency matrix B for $\tilde{\Gamma}$ yield the neighbourhood 1-(4060, 2304, 2304) design \mathcal{D}_{2304} . This is a self-dual symmetric design, and the binary row span of gives the code \mathcal{C}_{2304} .

The graphs, designs and codes

- From the rows of an adjacency matrix $B + I$ of the reflexive graph $\tilde{\Gamma}^R$ we get the self-dual symmetric 1-(4060, 2305, 2305) design \mathcal{D}_{2305} . We write \mathcal{C}_{2305} to denote the binary code of \mathcal{D}_{2305} .

Results

Lemma

Let G be the Rudvalis group Ru and \mathcal{D}_i and C_i where $i \in \{1755, 2305, 4059\}$ be the designs and binary codes constructed from the primitive rank-3 permutation action of G on the cosets of $2_{F_4(2)}$. Then

- (i) $\text{Aut}(\mathcal{D}_{1755}) = \text{Aut}(\mathcal{D}_{2305}) = Ru$ and \mathcal{D}_{1755} is the *unique point-primitive and flag-transitive symmetric design on 4060 points*.
- (ii) $C_{1755} = C_{2305} = V_{4060}(\mathbb{F}_2)$.
- (iii) $\text{Aut}(C_{1755}) = \text{Aut}(C_{2305}) = S_{4060}$.

Sketch of the proof

Proof: (i)

- The definition of Ω and \mathcal{B} emerges from Result 1.3, and from this it is clear that $G \subseteq \text{Aut}(\mathcal{D}_{1755})$.
- It follows from Result 1.3, and also from the Atlas [1, p.126] that G acts primitively on both Ω and \mathcal{B} of degree $|\Omega| = |\mathcal{B}| = 4060$, and the stabilizer of a vertex u (point) has exactly three orbits in Ω .
- G_u fixes setwise each of $\{u\}$, Δ_1 and $\Omega \setminus (\Delta_1 \cup \{u\}) = \Delta_2$ and these are all possible G_u -orbits.
- \mathcal{D}_{1755} is a point primitive, symmetric 1-design. It remains to show that $G = \text{Aut}(\mathcal{D}_{1755})$.
- Now $G \subseteq \text{Aut}(\mathcal{D}_{1755}) \subseteq S_{4060}$, so $\text{Aut}(\mathcal{D}_{1755})$ is a primitive permutation group on Ω of degree 4060. Moreover, $\text{Aut}(\mathcal{D}_{1755})_u$ must fix Δ_1 setwise, and hence $\text{Aut}(\mathcal{D}_{1755})_u$ also has orbits of lengths 1, 1755, and 2304 in Ω .

Sketch of the proof

- The only primitive group of degree 4060, such that $\text{Aut}(\mathcal{D}_{1755})_u$ can have orbit lengths 1, 1755, and 2304 is Ru , see [3, Theorem 18].
- $G = \text{Aut}(\mathcal{D}_{1755})$. Since $\mathcal{D}_{2305} = \tilde{\mathcal{D}}_{1755}$, we deduce that $\text{Aut}(\mathcal{D}_{2305}) = \text{Aut}(\mathcal{D}_{1755}) = \text{Ru}$.
- Recall that there is a unique class of maximal subgroups of Ru of type $2_{F_4(2)}$. Now, given a subgroup K in that class, its normalizer is twice bigger in Ru , meaning that there are exactly two subgroups $2_{F_4(2)}$ that contain K , and so we derive a contradiction.
- Thus, we conclude that there is a unique 1-(4060, 1755, 1755) symmetric design invariant under Ru , and since the block stabilizer acts transitively on the points of the block the claim on flag-transitivity holds.

The code of the graph Γ^R

Lemma

For Ru of degree 4060, the automorphism group of the graph Γ^R or design \mathcal{D}_{1756} is a non-abelian finite simple group of order 145926144000. Moreover this group is isomorphic to the simple sporadic group Ru .

Proof: This follows readily by computations with Magma. ■

Lemma

*The group Ru is the automorphism group of the $[4060, 29, 1756]_2$ code C_{1756} obtained from \mathcal{D}_{1756} . The code C_{1756} is **self-orthogonal doubly-even**. Its dual is a $[4060, 4031, 4]_2$ code. Moreover, $j \in C_{1756}$.*



The code of the graph $\tilde{\Gamma}$

Lemma

For R_u of degree 4060, the automorphism group of the design \mathcal{D}_{2304} is isomorphic to the group R_u .

Proof: Since $\mathcal{D}_{2304} = \tilde{\mathcal{D}}_{1756}$, we have $\text{Aut}(\mathcal{D}_{2304}) = \text{Aut}(\tilde{\mathcal{D}}_{1756}) = \text{Aut}(\mathcal{D}_{1756})$. Now the proof follows from Lemma 1.5. ■

Lemma

The group R_u is the automorphism group of \mathcal{C}_{2304} . The code \mathcal{C}_{2304} is *self-orthogonal doubly-even, with minimum weight 1792*. Its dual is a $[4060, 4032, 4]_2$. Moreover, *R_u acts irreducibly on \mathcal{C}_{2304} as an \mathbb{F}_2 -module, $\mathcal{C}_{2304} \subset \mathcal{C}_{1756}$, and $\text{Aut}(\mathcal{C}_{2304}) = R_u$.*



Sketch of the proof

Proof:

- Use the strong regularity of $\tilde{\Gamma}$ to show that the code C_{2304} is self-orthogonal.
- Notice first that C_{2304} is obtained from the strongly regular graph $\tilde{\Gamma}$ with parameters $(4060, 2304, 1328, 1280)$ and intersection matrix

$$\begin{bmatrix} 0 & 1 & 0 \\ 2304 & 1328 & 1280 \\ 0 & 975 & 1024 \end{bmatrix}.$$

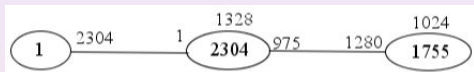
Sketch of the proof

- It can be seen from Figure 1 below that if we fix a vertex v in $\tilde{\Gamma}$ we can divide the remaining vertices into two sets, namely $\tilde{\Gamma}'$ of size 2304 and $\tilde{\Gamma}''$ of size 1755, with $\tilde{\Gamma}'$ being the set of vertices adjacent to v , and $\tilde{\Gamma}''$ the set of vertices non-adjacent to v .
- Now, from the second column of the above matrix we deduce that each vertex in $\tilde{\Gamma}'$ is adjacent to v and to 1328 other vertices in $\tilde{\Gamma}'$, thus to 975 vertices in $\tilde{\Gamma}''$ while from the third column shows that a vertex in $\tilde{\Gamma}''$ is adjacent to 1280 vertices in $\tilde{\Gamma}'$, and so to 1024 vertices in $\tilde{\Gamma}''$.

Sketch of the proof

- The structure of the graph and the orbit joins are summarized in the following diagram.

Figure : Number of joins between orbits of a stabilizer







- The valency 2304 ensures that generating codewords have length zero (mod 2) and the 1328 and the 1280 ensure that (i) any two generating codewords have an even number of non-zero entries in common, and (ii) that any two generating codewords are orthogonal to one another.
- Hence C_{2304} is self-orthogonal, and since all non-zero codewords have weights divisible by 4, it follows that C_{2304} is doubly-even.

Sketch of the proof



$$\begin{aligned}
 W_{C_{2304}} = 1 &+ 188500 x^{1792} + 4677120 x^{1952} \\
 &+ 38001600 x^{1984} + 95769600 x^{2016} \\
 &+ 95597775 x^{2048} + 33779200 x^{2080} \\
 &+ 417600 x^{2240} + 4060 x^{2304}.
 \end{aligned}$$

- Moreover, the blocks of \mathcal{D}_{2304} are of even size, so j meets evenly every vector of C_{2304} , so $j \in C_{2304}^\perp$. It can be deduced from [2, Section 3] that the 2-rank of $\tilde{\Gamma}$ is 28, and so the dimension of C_{2304} follows.
- If $\alpha \in \text{Aut}(C_{2304})$, then since $\alpha(j) = j$ and $C_{1756} = \langle C_{2304}, j \rangle$, we have $\alpha \in \text{Aut}(C_{1756})$. So that $\text{Aut}(C_{2304}) \subseteq \text{Aut}(C_{1756})$.
- Arguing similarly as Lemma 1.6 we show that $\text{Aut}(C_{2304}) = \text{Ru}$. ■

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Correction to: “Codes, designs and graphs from the Janko groups J_1 and J_2 ” [*J. Combin. Math. Combin. Comput.* **40** (2002), 143–159],

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